

1.12 Rules of inference with propositions

Rules of inference known to be valid arguments:

Rule of Inference	Name
$\begin{array}{l} p \\ p \rightarrow q \\ \hline \therefore q \end{array}$	Modus ponens
$\begin{array}{l} \neg q \\ p \rightarrow q \\ \hline \therefore \neg p \end{array}$	Modus tollens
$\begin{array}{l} p \\ \hline \therefore p \vee q \end{array}$	Addition
$\begin{array}{l} p \wedge q \\ \hline \therefore p \end{array}$	Simplification
$\begin{array}{l} p \\ q \\ \hline \therefore p \wedge q \end{array}$	Conjunction
$\begin{array}{l} p \rightarrow q \\ q \rightarrow r \\ \hline \therefore p \rightarrow r \end{array}$	hypothetical syllogism
$\begin{array}{l} p \vee q \\ \neg p \\ \hline \therefore q \end{array}$	Disjunctive syllogism
$\begin{array}{l} p \vee q \\ \neg p \vee r \\ \hline \therefore q \vee r \end{array}$	Resolution

Example

$$\begin{array}{l} \text{Argument : } (r \vee w) \rightarrow c \\ \quad \quad \quad \neg c \\ \hline \quad \quad \quad \therefore \neg w \end{array}$$

Proof:

1.	$(r \vee w) \rightarrow c$	hypothesis
2.	$\neg c$	hypothesis
3.	$\neg (r \vee w)$	Modus tollens 1,2
4.	$\neg r \wedge \neg w$	De Morgans Law, 3
5.	$\neg w \wedge \neg r$	Commutative Law, 4
6.	$\neg w$	Simplification, 5

1.13 Rules of inference with quantifiers

- A value that can be plugged in for variable is called an element of the domain.
 - An arbitrary element of the domain has no special properties other than those shared by all elements of the domain.
 - A particular element of the domain may have properties that are not shared by all the elements of the domain.
 - * if D is $\mathbb{R}$, then 3 is a particular element because it is odd.
- Any element defined in a hypothesis is particular.

• The rules of existential instantiation and universal instantiation replace a quantified variable with an element of the domain.

• Universal Instantiation says that any element in the domain can substitute for a universally quantified variable.

• Existential instantiation introduces a particular element to replace an existentially quantified variable.

* A new element with a new name must be used for each use of existential instantiation.

• The rules of existential generalization and universal generalization replace an element of the domain with a quantified variable.

• Existential generalization says that an element of the domain can be replaced by an existentially quantified variable

• The rules ONLY APPLY TO NON-NESTED QUANTIFIERS

Rule of Inference	Name
C is an element (arbitrary or particular) $\forall x P(x)$ $\therefore P(c)$	Universal instantiation
C is an arbitrary element $P(c)$ $\therefore \forall x P(x)$	Universal generalization
$\exists x P(x)$ $\therefore (C \text{ is a particular element}) \wedge P(c)$	Existential instantiation
C is an element (arbitrary or particular) $P(c)$ $\therefore \exists x P(x)$	Existential generalization

2.1 Mathematical definitions

- An integer x is even if there is an integer k such that $x = 2k$
- An integer x is odd if there is an integer k such that $x = 2k + 1$
- The parity of a number is whether the number is even or odd.
 - Two numbers have same parity if they are both even or both odd
 - Two numbers have opposite parity if one is even and one is odd
- A number r is rational if there exist integers x and y such that $y \neq 0$ and $r = \frac{x}{y}$.
- Divides
 - An integer x divides an integer y if and only if $x \neq 0$ and $y = kx$, for some integer k
 - The fact that x divides y is denoted $x \mid y$. If x does not divide y , the notation is $x \nmid y$
 - If x divides y , then y is a multiple of x , and x is a factor or divisor of y .
- An integer n is prime i.f.f. $n > 1$, and the only positive integers that divide n are 1 and n .
- An integer n is composite i.f.f. $n > 1$, and there is an integer m such that $1 < m < n$ and m divides n
- A real number x is positive i.f.f. $x > 0$
- A real number x is negative i.f.f. $x < 0$
- A real number x is nonnegative i.f.f. $x \geq 0$
- A real number x is nonpositive i.f.f. $x \leq 0$